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# **Pushing XPath Accelerator to its Limits**

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## Overview

**motivation**  
...processing large XML documents

**baseX & iDEx**  
...our two prototypes

**PerfidiX**  
...a benchmark framework

**evaluation**  
...performance results

**lessons learned**  
...what we will do next



# Motivation

## Observation

- sizes of XML instances are continuously growing:
  - Library Data, U Konstanz: 2 GB
  - DBLP: 300 MB
  - Wikipedia: 5 GB up to 500 GB
  - Log files > 10 GB ...

## Fact

- XML processors needed to handle these documents
- current XML processors usually fail:
  - by design (on-the-fly parsing, 2GB limit, indexing overhead, ...)
  - by technical limits (main memory barrier, swapping, ...)



# Motivation

## MonetDB/XQuery

- based on the Pathfinder project, developed in Konstanz
- XPath Accelerator: relational XML encoding
- StairCase Join: very efficient path traversal
- Loop-Lifting: linear execution of nested loops

## Identified Bottlenecks (Challenges...)

- main memory limitation
- no content/value indexes



## Motivation

### Two Approaches

#### BaseX

shrink main memory representation

- pure main memory processing
- compressed representation of XPath Accelerator encoding
- introduction of an inherent value index

#### Idefix

optimize disk layout

- persistent native XML storage ←
- constant scalability ←
- logarithmic updateability ←





# BaseX – Memory Architecture

## Node Table Representation

Pre	Par	Tag	Content	Kind	AttName	AttVal	Parent (32 bit)	Kind/Token (1/31 bit)	Attributes (10/22 bit)
1	0	db		elem			...0000	0....0000	nil
2	1	address		elem	id	add0	...0001	0....0001	0000...0000
3	2	name	Hack Hacklinson	elem	title	Prof.	...0010	0....0010	0001...0001
4	3			text			...0011	1....0000	nil
5	2	street		elem			...0010	0....0011	nil
6	3		Alley Road 43	text			...0011	1....0001	nil
7	2	city		elem			...0010	0....0100	nil
8	3		Chicago, IL 60611	text			...0011	1....0010	nil
9	1	address		elem	id	add1	...0010	0....0001	0000...0010
10	9	name		elem			...0011	0....0010	nil
11	10		Jack Johnson	text			...0001	0....0001	0000...0010
12	9	street		elem			...0010	0....0010	nil
13	10		Pick St. 43	text			...	...	...
14	9	city		elem					
15	10		Phoenix, AZ 85043	text					

ID	Tag	ID	Text	ID	AttName	ID	AttValue
0000	db	0000	Hack Hacklinson	0000	id	0000	add0
0001	address	0001	Alley Road 43	0001	title	0001	Prof.
0010	name	0010	Chicago, IL 60611			0010	add1
0011	street	0011	Jack Johnson				
0100	city	0100	Pick St. 43				
0101		0101	Phoenix, AZ 85043				

⇒ index storage

⇒ numeric references



# BaseX – Querying

## Value Indexing

- Text and AttributeValue indexes are extended by references to Pre values ( $\Rightarrow$  inverted index)
- small memory overhead (12 – 18%)

## Query Optimization:

- predicates are evaluated first (selection pushdown)
- internal `index` axis and `cs()` kind test are added for predicate evaluation
- queries are inverted & rewritten

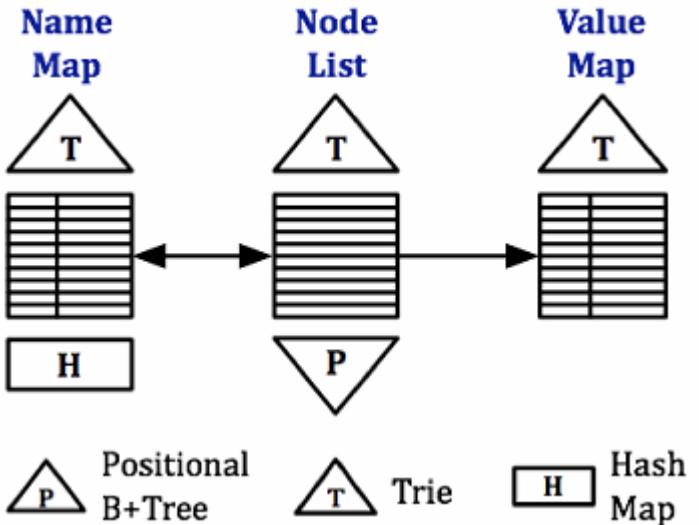
## Example:

`/db/address[@id = "add0"]/name`  
 $\Rightarrow$  `index::node()[@id = "add0"]/parent::address[parent::db/parent::cs()]/child::name`

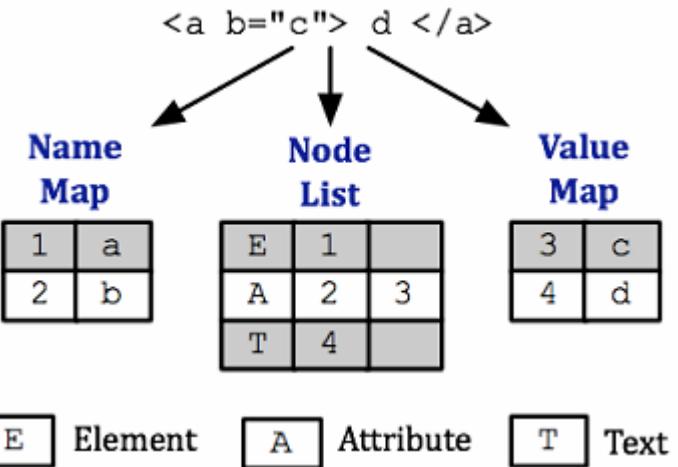


# Idefix – Data Structures

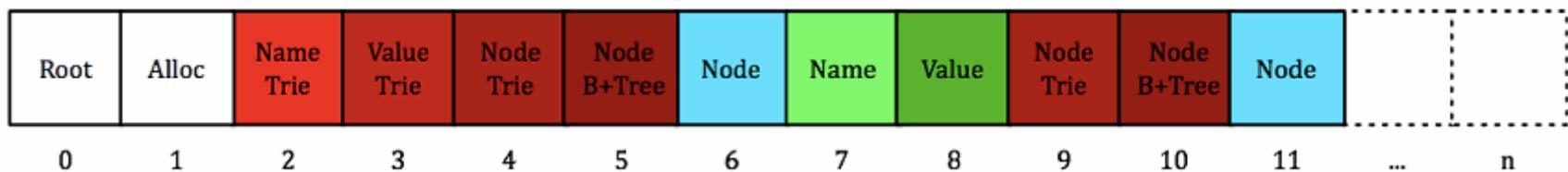
## Concept



## Shredding



## Block Storage





# Perfidix – Java Benchmarking Framework

## Task

- automate tedious manual benchmarking tasks
  - generic
  - à la JUnit
  - integration (Eclipse, Ant, ...)

## Output

- console or XML per benchmark ( $n$  runs)
  - minimum, maximum, average, standard deviation, confidence interval 95

## Discussion

- Java memory management
- benchmark history



# Perfidix – Java Benchmarking Framework (cont.)

## Code Example

```
public class DemoBench extends Benchmarkable {  
    public DemoBench() {...}                      // one-time initialization  
    public setUp() {...}                          // per-run & method preparation  
    public tearDown() {...}                     // per-run & method cleanup  
    public benchFoo() {...}                      // method Foo to bench  
    public benchBar() {...}                      // method Bar to bench  
}
```

## Output Example

-	unit	sum	min	max	avg	stddev	conf95	runs
benchFoo	ns	8023000	19000	3822000	80230.00	376167.54	[6501.16,153958.84]	100
benchBar	ns	3951000	15000	778000	39510.00	74585.05	[24891.33,54128.67]	100
TOTAL	ns	11974000	3951000	8023000	5987000.00	2036000.00	[3165247.96,8808752.04]	



# Evaluation

## Systems

- MonetDB & BaseX
  - main memory based processing
  - similar data structures
- X-Hive & Idefix
  - persistent disk storage
  - comparable scalability

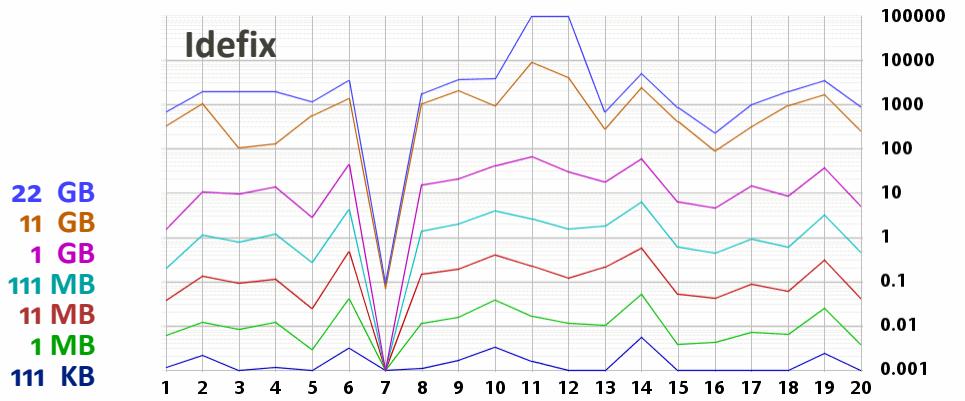
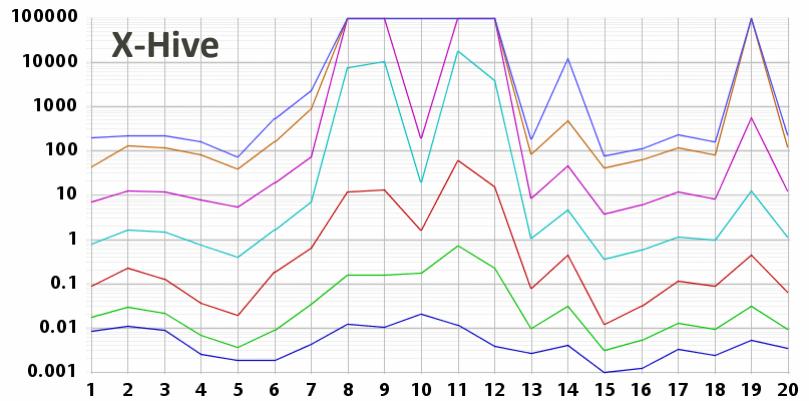
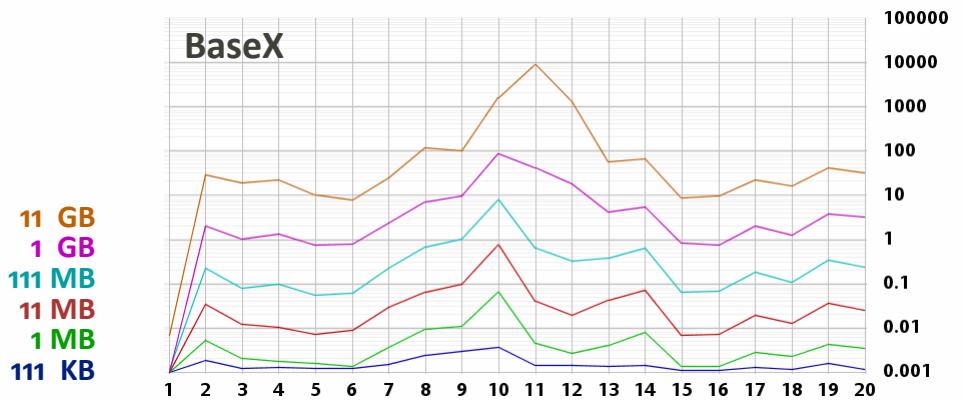
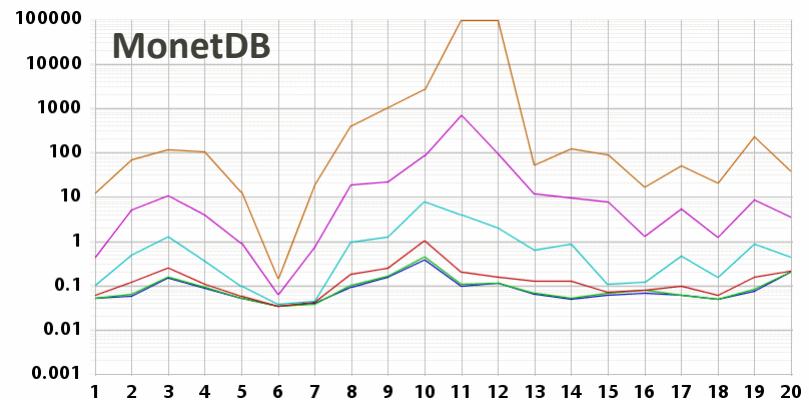
## Benchmark Queries

- XMark, 110 KB – 22 GB
- six value-oriented DBLP Queries, 300 MB



## Evaluation – Scalability

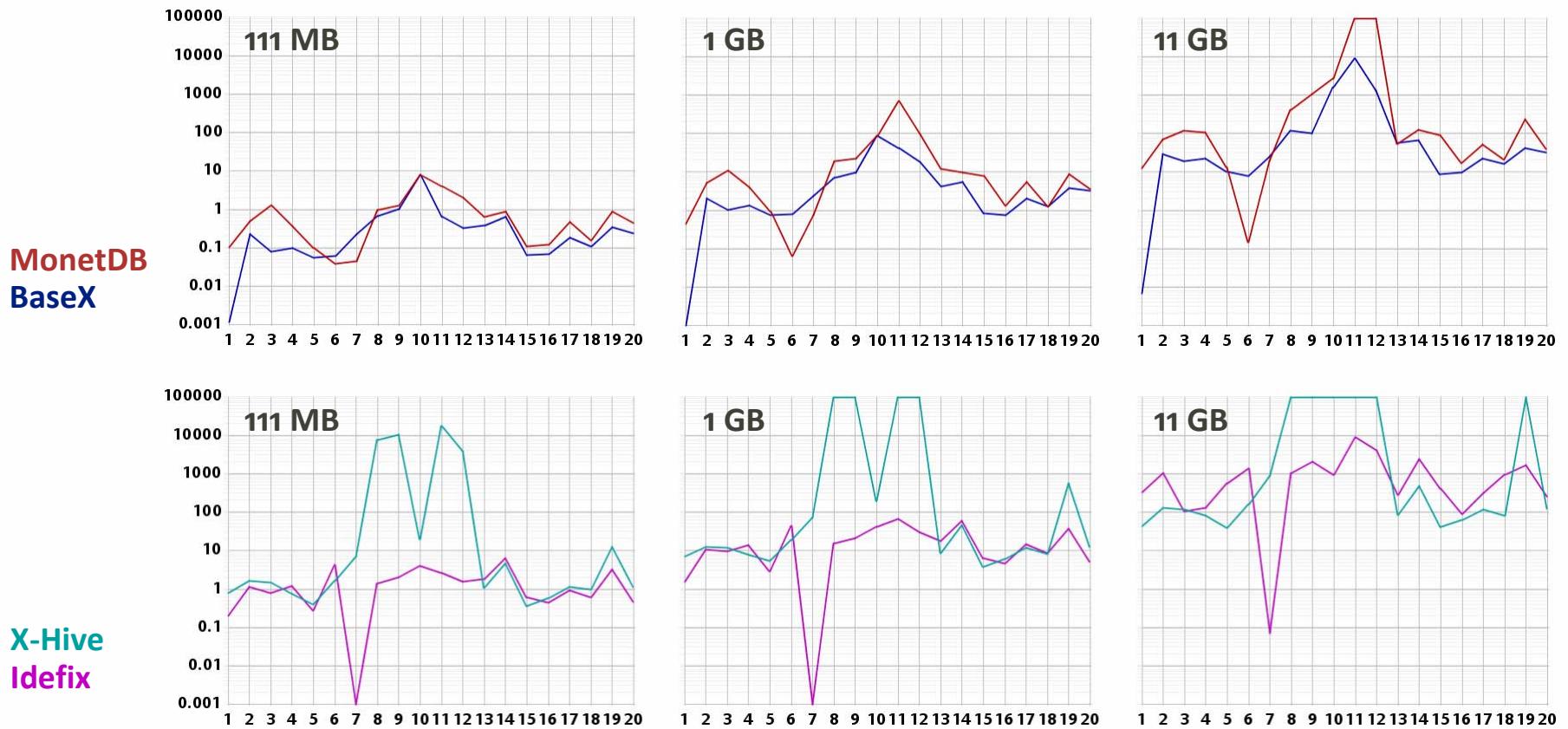
XMark queries (x-axis ⇒ number of query, y-axis ⇒ execution time in sec.)





## Evaluation – XMark

XMark queries (x-axis ⇒ number of query, y-axis ⇒ execution time in sec.)





## Evaluation – DBLP

DBLP queries (x-axis: number of query, y-axis: execution time in sec.)

contains() function:

```
[1] /dblp/*[contains(title, 'XPath')]
```

range query:

```
[2] /dblp/*[year/text() < 1940]/title
```

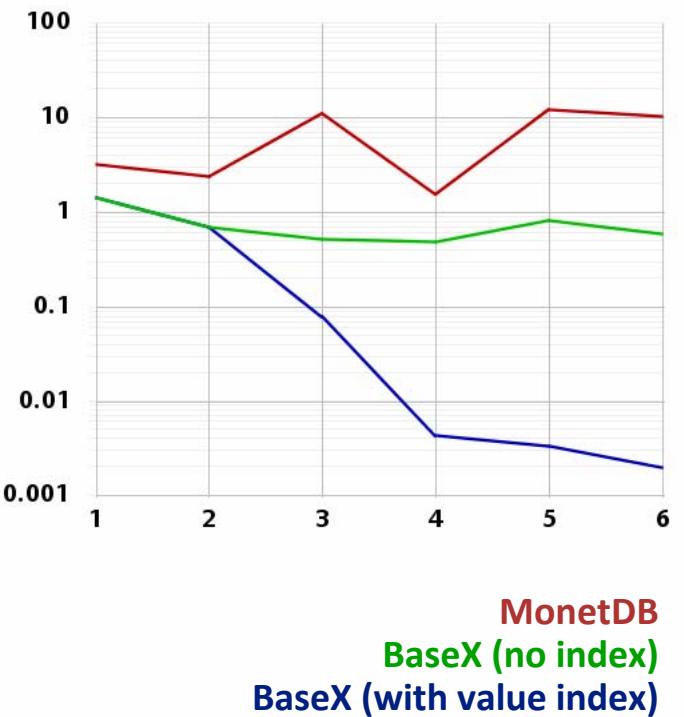
exact predicate match:

```
[3] /dblp//inproceedings[contains(@key, '/edbt/')]  
    [year/text() = 2004]
```

```
[4] /dblp/article[author/text() = 'Alan M. Turing']
```

```
[5] //inproceedings[author/text() = 'Jim Gray']/title
```

```
[6] //article[author/text() = 'Donald D. Chamberlin']  
    [contains(title, 'XQuery')]
```





# Lessons Learned

- hard-coded queries might blur evaluation results
- comparison troublesome with different systems
  - granularity of measurements (shredding, compilation, serialization, ...)
  - impact of different system components (storage, query)
  - availability of different features (updates, complete query implementation)
- handling of serialization output
- assure correctness of large results
- many factors to measure:
  - CPU load
  - memory I/O
  - disk I/O
  - memory consumption



# Future Work

## Merge BaseX & Idefix

- comprehensive support for value-based queries
- full text queries, including scoring algorithms (like SRA/INEX)
- optimize XML table compression
- optimize disk layout (hybrid, networked, and holographic storage)
- write Pathfinder plugin to support XQuery
- complete update implementation

## Benchmarking

- use of virtual machines for benchmark reproducibility
- specify benchmark for XML updates
- application benchmark for eMail storage